CSE 3302/5307 Programming Language Concepts

Homework3 - Fall 2023

Due Date: Sep.16, 2023, 11:59p.m. Central Time

Problem1 - 20%

Evaluate $(\lambda x. ((\lambda y. x + z + 3) 3)5)$ using call-by-value and call-by-name. Show the complete steps of evaluation.

Problem2 - 20%

Given $Y = \lambda f(\lambda x. f(xx))(\lambda x. f(xx))$, show that g(Yg) = Yg by performing call-by-name evaluation on Yg.

Problem3 - 30%

Prove by induction: If $FV(e_1) = \emptyset$ and $e_1 \to e_2$, then $FV(e_2) = \emptyset$.

- Given the following definitions:
 - 1. Rules of free variables

$$\frac{FV(e_1) = S_1 \quad FV(e_2) = S_2}{FV(e_1 e_2) = S_1 \cup S_2} \qquad \frac{FV(e) = S}{FV(x.e) = S - \{x\}}$$

2. Judgment form: **define** $e_1 \rightarrow e_2$

$$\frac{e_1 \to e_1'}{(\lambda x.e) \ v \to e[v/x]} \qquad \frac{e_1 \to e_1'}{e_1 \ e_2 \to e_1' \ e_2} \qquad \frac{e_2 \to e_2'}{v \ e_2 \to v \ e_2'}$$

• And given this lemma:

Lemma 1.
$$FV(e_1[e_2/x]) \subseteq (FV(e_1) - \{x\}) \cup FV(e_2)$$

By induction on derivation of $e_1 \to e_2$ 1. Case $\frac{1}{(\lambda x.e) \ v \to e[v/x]}$ Need to Prove:

2. Case $\frac{e_1 \rightarrow e_1'}{e_1 \ e_2 \rightarrow e_1'}$ Need to Prove:

3. Case $\frac{e_2 \rightarrow e_2'}{v \ e_2 \rightarrow v \ e_2'}$ Need to Prove:

Problem4 - 30%

Church numerals use lambdas to create a representation of numbers. They can represent natural numbers 0, 1, 2, ..., as follows:

$$\mathbf{0} = \lambda f.\lambda x. \ x$$

$$\mathbf{1} = \lambda f.\lambda x. \ f \ x$$

$$\mathbf{2} = \lambda f.\lambda x. \ f \ (f \ x)$$

$$\mathbf{3} = \lambda f.\lambda x. \ f \ (f \ (f \ x))$$
...
$$\mathbf{n} = \lambda f.\lambda x. \ f^n \ x$$

Church numerals takes two parameters f and x. Church numerals n means apply f to x n times. You can read more about church numerals on the internet.

(a) Define addition in λ calculus, and then show the evaluation of 3+2.

(b) Define multiplication in λ calculus (Hint: you can use definition of addition), and then show the evaluation of 3×2 .